

# Smartcard-Login into Gemstone

An exercise in applied security

by

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# Disclaimer

-----BEGIN PGP SIGNED MESSAGE-----

Hash: SHA1

This Material is a copy of the presentation at the ESUG 2003, Bleed, Slovenia.

It has been developed by Alfred Wullschleger,  
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-----BEGIN PGP SIGNATURE-----

Version: PGP 8.0

iQA/AwUBP1MIRIBkqNDRsW06EQINYACg63d2yK0GCdIz0f21QcfWupyx1fwAoO4Q  
px5moaJuaAkPTCCgtxq5+iNK  
=bMoJ

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# Talk and Technical Discussion

- Two Parts
  - first Part (this Talk): Presentation of the ideas
  - second Part: Technical discussion
    - code examples
    - demonstrations
    - details of the implementation
      - as requested by the audience

# The Author

- Smalltalker since 1992, never ending enthusiasm
- Project OVID at FIDES Informatik (1992-1999)
- Leading the project OASE at SNB since 1999
  - Financial statistics from Swiss banks and Swiss companies
  - Fully object oriented application using VW3.0 and Gemstone as development base

# Motivation

- Gemstone Login uses username and password
  - Both are sent over the network in clear text
  - passwords are typically „trivial“ by User requirement?!
- Smartcards (SC) have many advantages
  - Very high security
    - User must have the card and a password for the card
  - Even relatively trivial passwords are acceptable since the SC is completely disabled after 3 trials

# Smartcards at SNB

- Smartcards will be introduced in the Bank during 2004
- Goal is Single Signon
  - OASE is one of the first applications which today are ready for Smartcard logon

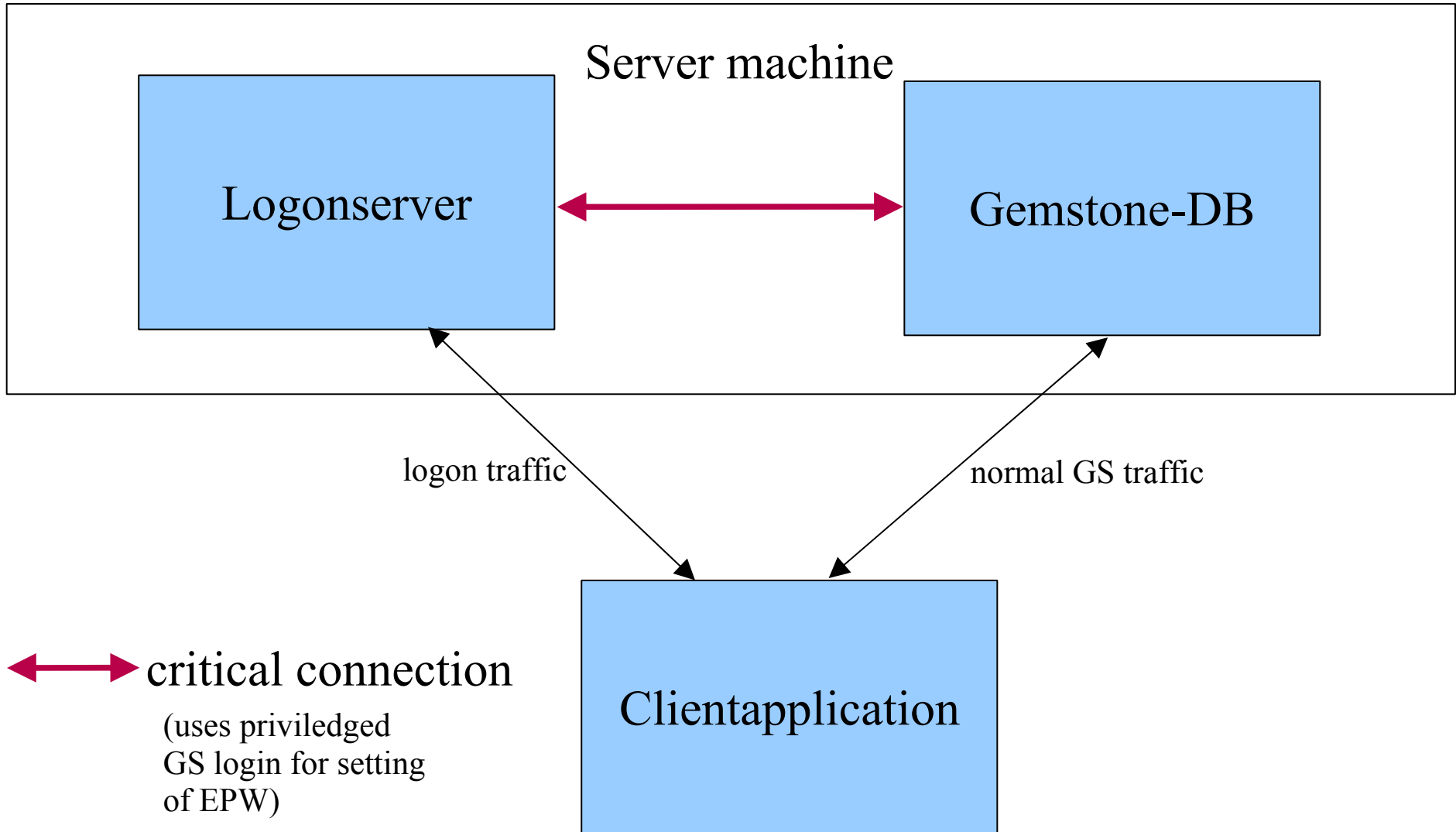
# Smartcard Login

How to implement?

# Gemstone Prerequisites

- Logon uses username and password
  - Password size 1024 byte maximum
  - Password expiration can be set to exactly one logon
- So, we can use a „complicated“ One-Time-Password, we call it EPW (from german)
  - format can be choosen appropriate
- Smartcard can be combined with GS:
  - use a Logonserver

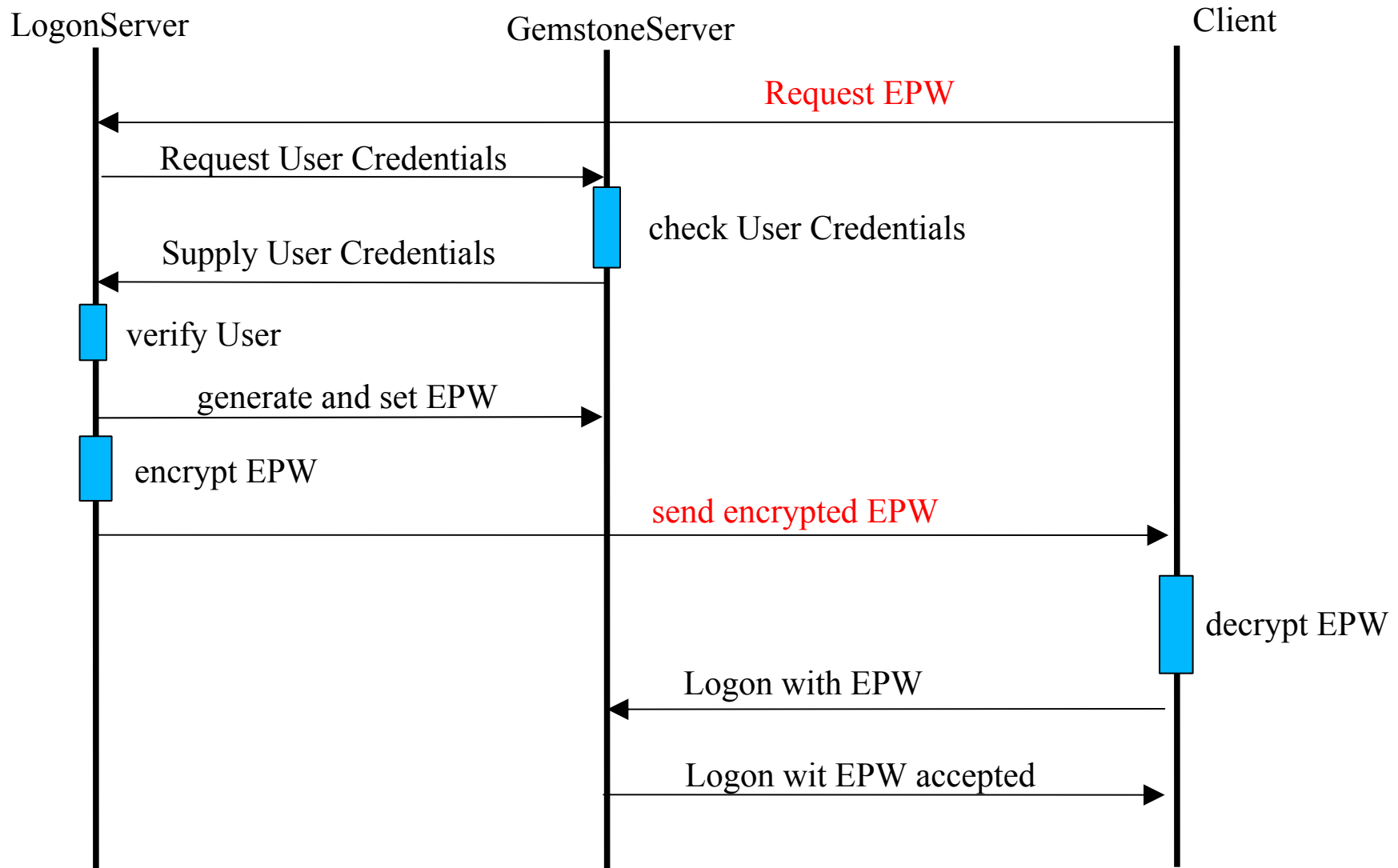




# Logonserver (LS)

- LS runs on the same machine as the GS server(s)
  - this guarantees the exchange of critical data with GS without network traffic
  - handles all GS servers on the same machine
  - handles all client logon requests for these servers
  - is implemented as a headless VW3.0 application
    - so it runs wherever we want (AIX, Win2000, Linux, WinXP...)
    - runs as TCP/IP-Server

# LS basic operation



red = encrypted

# LS basic operation

- the client requests his EPW from LS
- LS verifies with GS, that user is ok
  - (i.e. has credentials and a user account)
- LS generates EPW and sets it in GS for one use
- LS encrypts the EPW and sends it to the client
- the client decrypts the EPW and makes normal GS login with his username and the EPW

# LS has two concurrent modes

- **SC-RSA (Smartcard-RSA-mode)**
  - The user uses a Smartcard with RSA-signature keypair
- **PW-DH (Password-Diffie-Hellman-mode)**
  - The user has no Smartcard
    - a forgotten Smartcard should not prevent the user from logging into GS
    - this mode must be activated in GS for each user individually by administrators, if in SC-RSA-mode before

# We need a secure data exchange Client $\leftrightarrow$ LS

- Exchange of secret data through encryption
  - we use Blowfish block cipher
- Each session Client-LS must have a session key
  - can be generated by using Diffie-Hellman (DH) technique
  - how does DH work?

# Diffie-Hellman part1

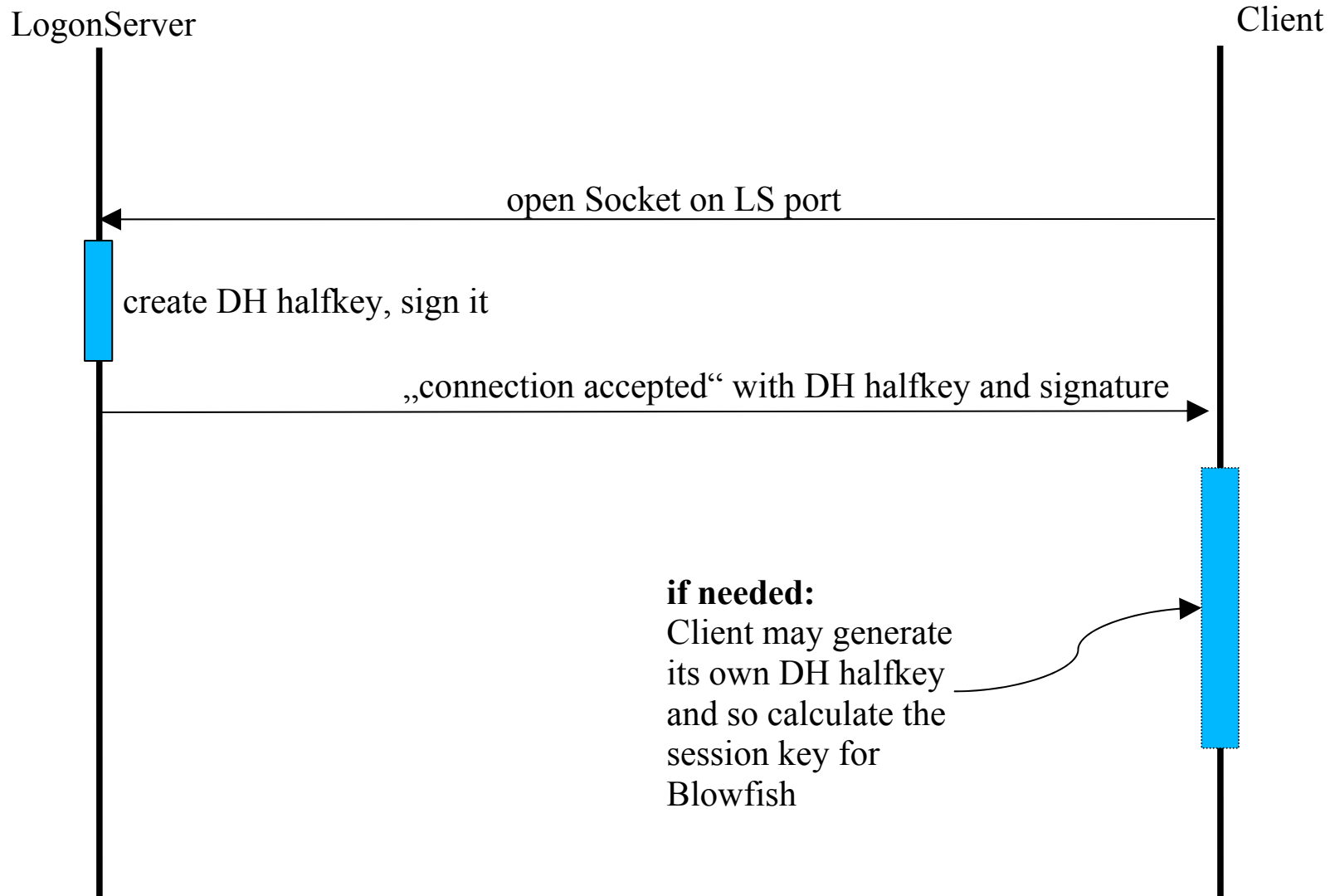
- creation of one secret for two parties
- based on the cyclic group  $\mathbf{Z}_p = \{1,2,3\dots p\}$ , where  $p$  is prime.
  - this is equivalent to multiplication of natural numbers modulo  $p$
  - $p$  size large, typically  $> 256$  bits
- use a common base  $b \in \mathbf{Z}_p$ , e.g. 65537
- each party selects a secret random number  $r_j > 1$ , typically large.  $j=1,2$

# Diffie-Hellman part2

- each party calculates  $\text{pub}_j = (b^{**} r_j) \bmod p$ ,  
so,  $\text{pub}_j \in \mathbf{Z}_p$
- both parties publicly exchange their  $\text{pub}_j$ 
  - we call  $\text{pub}_j$  the DH-Halfkey of party  $j$
- each party calculates  $\text{key} = (\text{pub}_j)^{**} r_i \bmod p$ 
  - ( $j = \text{other}, i = \text{self}$ )
- Since  $(b^{**} r_i)^{**} r_j == (b^{**} r_j)^{**} r_i$ , both parties get the same key!



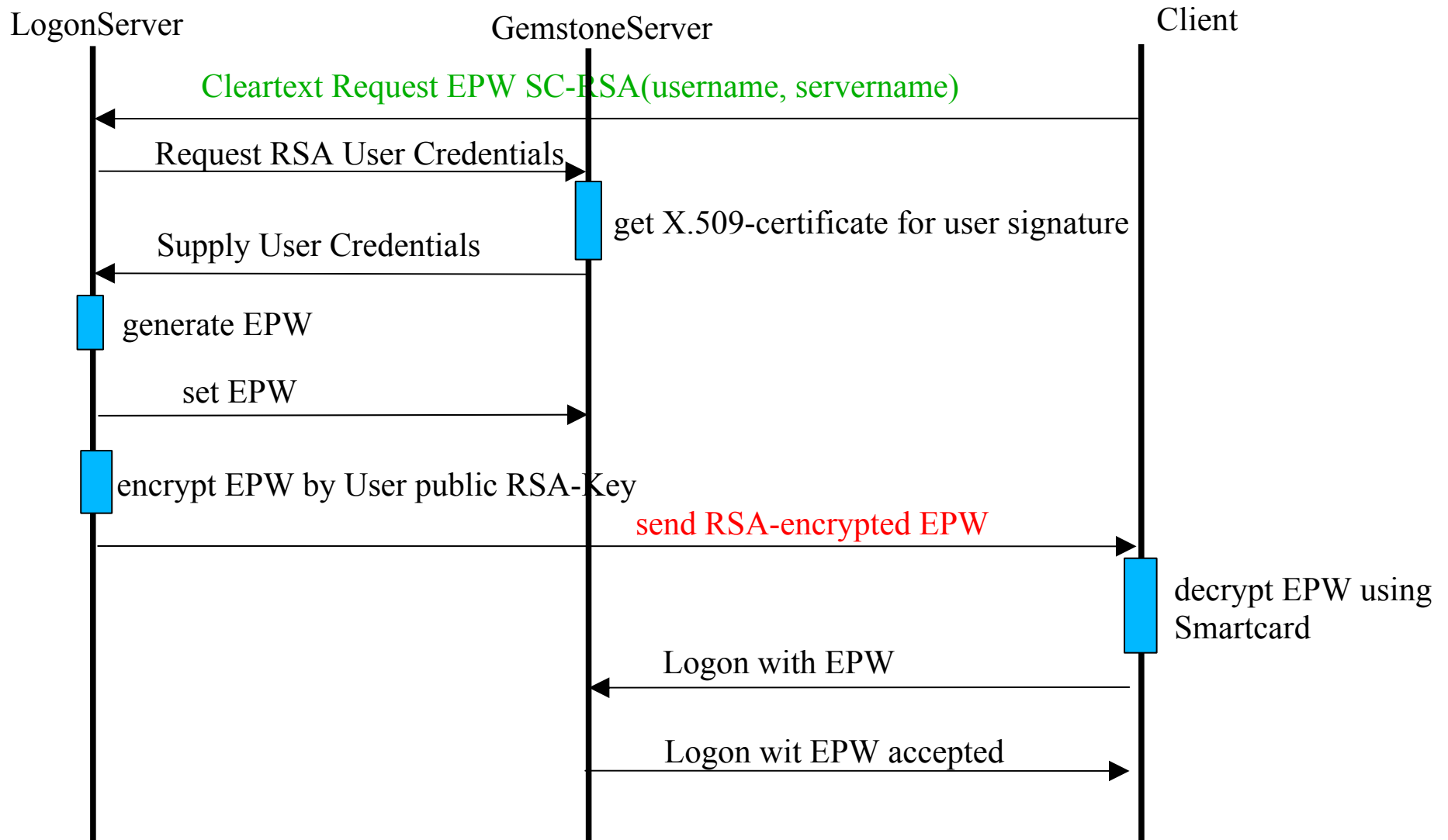
# LS connection start



# LS connection start (both modes)

- As soon as a connection to LS is established
  - LS creates a Diffie-Hellman halfkey for generation of a session key
  - LS signs this halfkey, so that a user can verify, that he is connected to a legal LS
  - LS sends this information as a „connection accepted“ message to the client
  - this is done irrespective of the mode wished by the client

# SC-RSA Overview



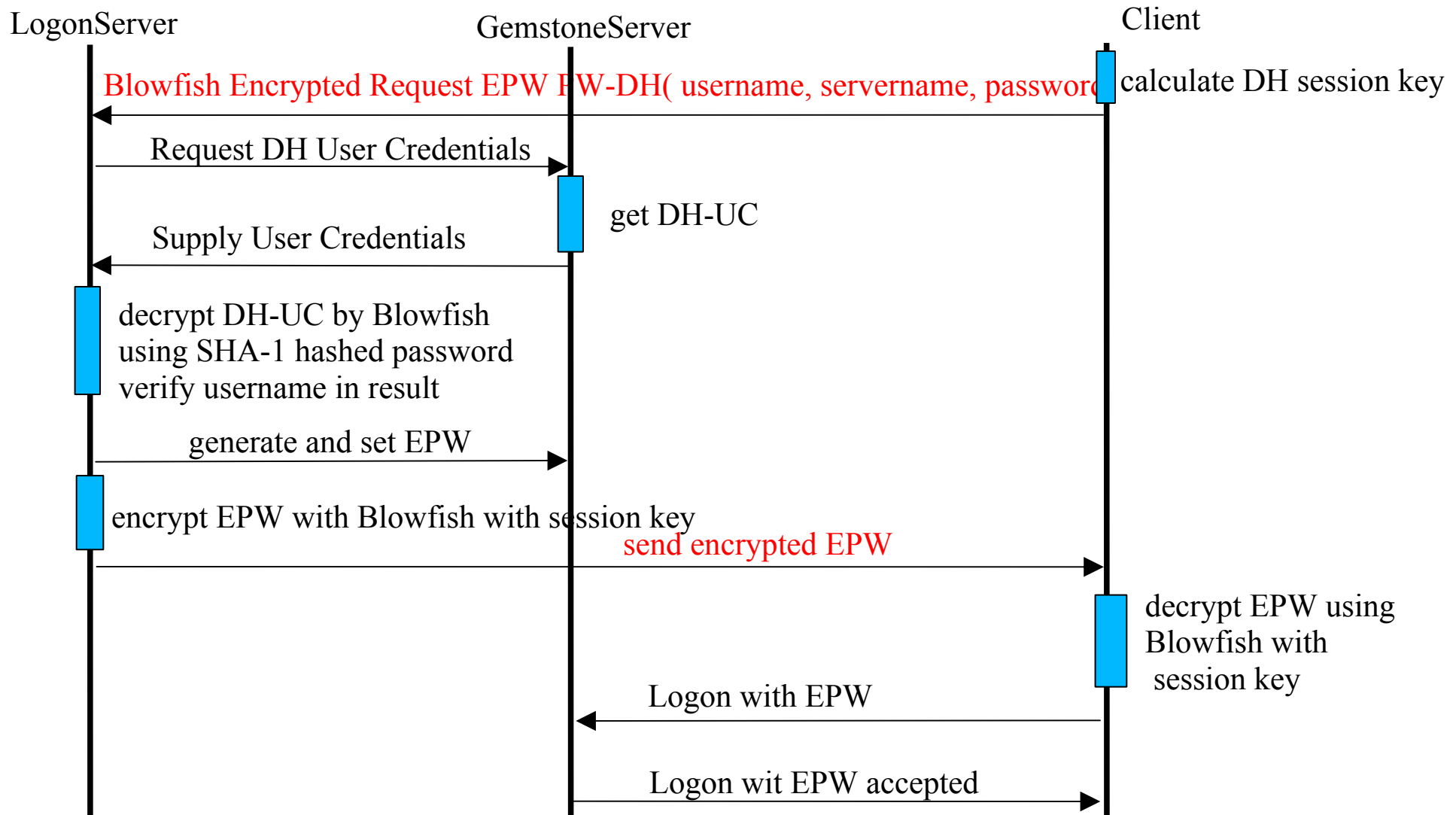
# SC-RSA part1

- User requests EWP in SC-RSA-mode
  - parameters: GS servername, GS username
  - sent in clear text to LS
- LS requests X.509-Certificate for the user from GS
  - LS generates EPW and sets it in GS
  - LS RSA-encrypts EPW using signaturekey from the X.509-Certificate

# SC-RSA part2

- LS sends encrypted EPW to client
  - if the user has the corresponding Smartcard, he can decrypt the EPW, if not, the information from LS is useless
- user decrypts EPW with his Smartcard
  - makes login in GS, using GS username and EPW

# PW-DH Overview



# PW-DH part1

- client requests EWP in PW-DH-mode
  - parameters: GS servername, GS username, password
  - client generates halfkey for DH-Keyexchange
  - creates with the halfkey from LS the session key for blowfish encryption
  - sends EPW request encrypted to LS together with his halfkey
- LS decrypts EPWrequest
  - by that gets servername, username and password

# PW-DH part2

- LS gets encrypted PW-DH-credentials for username from GS
  - GS stores PW-DH-credentials as blowfish-encrypted username plus randomdata
  - password is hashed with SHA-1 to get a blowfish key
    - Watch out: hashing does not help if password is trivial
  - These data are generated by interaction of an administrator with the user
    - when RSA-SC is used, PW-DH-credentials are erased



# PW-DH part3

- LS decrypts PW-DH-credentials using the password, if ok:
  - LS generates EPW and sets it in GS
  - LS blowfish-encrypts EPW using the sessionkey
- client receives and decrypts EPW
  - makes login in GS, using GS username and EPW

# Implementation issues

Encryption,  
Formats of EPW etc.

# RSA basics

- use two secret primes  $p$  and  $q$ :
  - get  $n = p * q$ , a large number (1024 bit)
  - select a number  $e$  with  $\text{gcd}(e, (p-1) * (q-1)) = 1$
  - $e$  typically 65537 (or 17 or 3)
    - $n$  and  $e$  form the public key for encryption
  - calculate  $d$  so that  $ed \equiv 1 \pmod{(p-1) * (q-1)}$ 
    - $n$  and  $d$  form the secret decryption key

# RSA implementation

- Security is based on factorization being a hard problem
- RSA is fully implemented in VW7.\*
  - we have ported it to VW3.0
    - was very easy

# Blowfish

- Block cipher defined by Bruce Schneier, 1994
  - encrypts 64-bit-Blocks
- uses keylength between 128 and 448 bit
  - we use 256 bit keys
- we use CBC-Mode
  - each block is XORed with the preceeding encryptedBlock, first block with an InitialVector

# Sizes

- RSA-keys on SC and signing key of LS: 1024bit
- Diffie-Hellman parameters
  - p: currently 256 bit
  - base = 65537
- EPW
  - 20-byte random ByteArray converted to LargePositiveInteger
    - giving a passwordstring of 48 to 49 digit characters

# XML for all requests/responses

- XML as exchange format between Client and LS
  - To avoid escape sequences like `&#x0027;` etc., we are using integer formats throughout all exchanges which use encryption
  - The `ByteArrays` are converted using „`asLargePositiveInteger`“
    - this is compatible with PKCS#1-specification of conversion between bytes and integers

# LS Answer on Socket accept:

```
- <XmlResponse>  
  <Validation>true</Validation>  
  <DhCode>1588342583341664246536080884883637711967782785292526  
  5980146286745390962263576</DhCode>  
  <ServerDHSignatur>1837090060268423857924167053495562451884751  
  288762027596014818947508697902699959279787764051953251210518  
  835088865241327868618698753687189866681129504242203066018314  
  921615856755804950286871069415444994298422094470968257289272  
  110594464664365849721683826653601332531687887166125489512675  
  5085488495158339417264534</ServerDHSignatur>  
</XmlResponse>
```

- User can verify the Signature of the LS



# SC-RSA basic format

- SC-RSA-request from client to LS is unencrypted:
  - `<RSAEPWRequest>`  
`<Server>lightning</Server><User>rog</User>`  
`</RSAEPWRequest>`
- SC-RSA-response from LS to client:
  - `<EinmalPasswort>254036481566762758850660080224084753875223586`  
`838390721369209245892308754789435119113869732769585450706139`  
`953046640274448442599416470993396942547126224650088959281761`  
`150591890399102073665906605091766315849615345087595361055920`  
`327412405609537797569746156102333670979136662194517047330045`  
`30461339320939262107515</EinmalPasswort>`

# SC-RSA improved

- we could improve the protocol by requesting, that the client has to sign the LS halfkey
  - so, we avoid setting EPW for a user, who has no corresponding X.509-certificate on GS for the SC
    - has more computing overhead
  - could avoid denial-of-service-attacks
    - disable user, if repeatedly illegally requesting an EPW
- this is implemented for ticket requests (see later)

# PW-DH request

- Internal PW-DH-Request:
  - `<XmlEPWRequest>`  
`<Server>lightning</Server>`  
`<User>awu</User><Password>xxxxxx</Password></XmlEPWRequest>`
- encrypted PW-DH-Request:
  - `<XmlEPWRequest>`  
`<DhCode>2815747482711583432974205332271263279346890877810142`  
`0710560243267932546510264</DhCode>`  
`<EinmalPasswortRequest>421701460956386515843309546275011481111`  
`952142904853134983891656694841409986396784019003171859680754`  
`920531201111390893851148664959465630186555935507756945874285`  
`836772314488857119582594456372386721952350160342234743728624`  
`71711709551746300760031120608509510132548629714428</EinmalPass`  
`wortRequest>`  
`</XmlEPWRequest>`

# PW-DH response

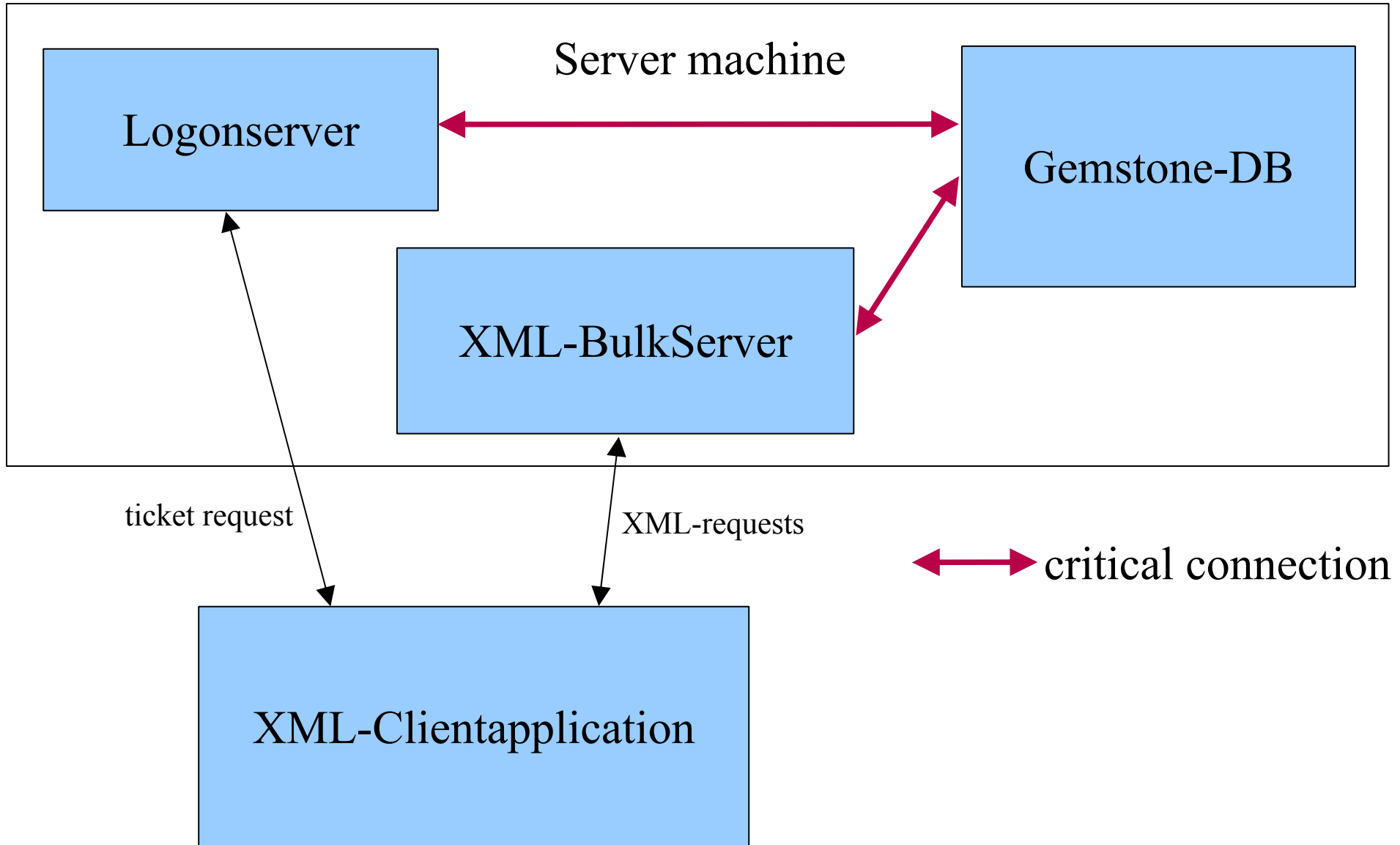
- Encrypted Response from LS:
  - `<EinmalPasswort>391016871021610782724341211645684570668712594  
970979368692619795917841928462118545730795848292112003218739  
3806816930803570651576119091883739061802251407795  
</EinmalPasswort>`

# Other Functions of LS

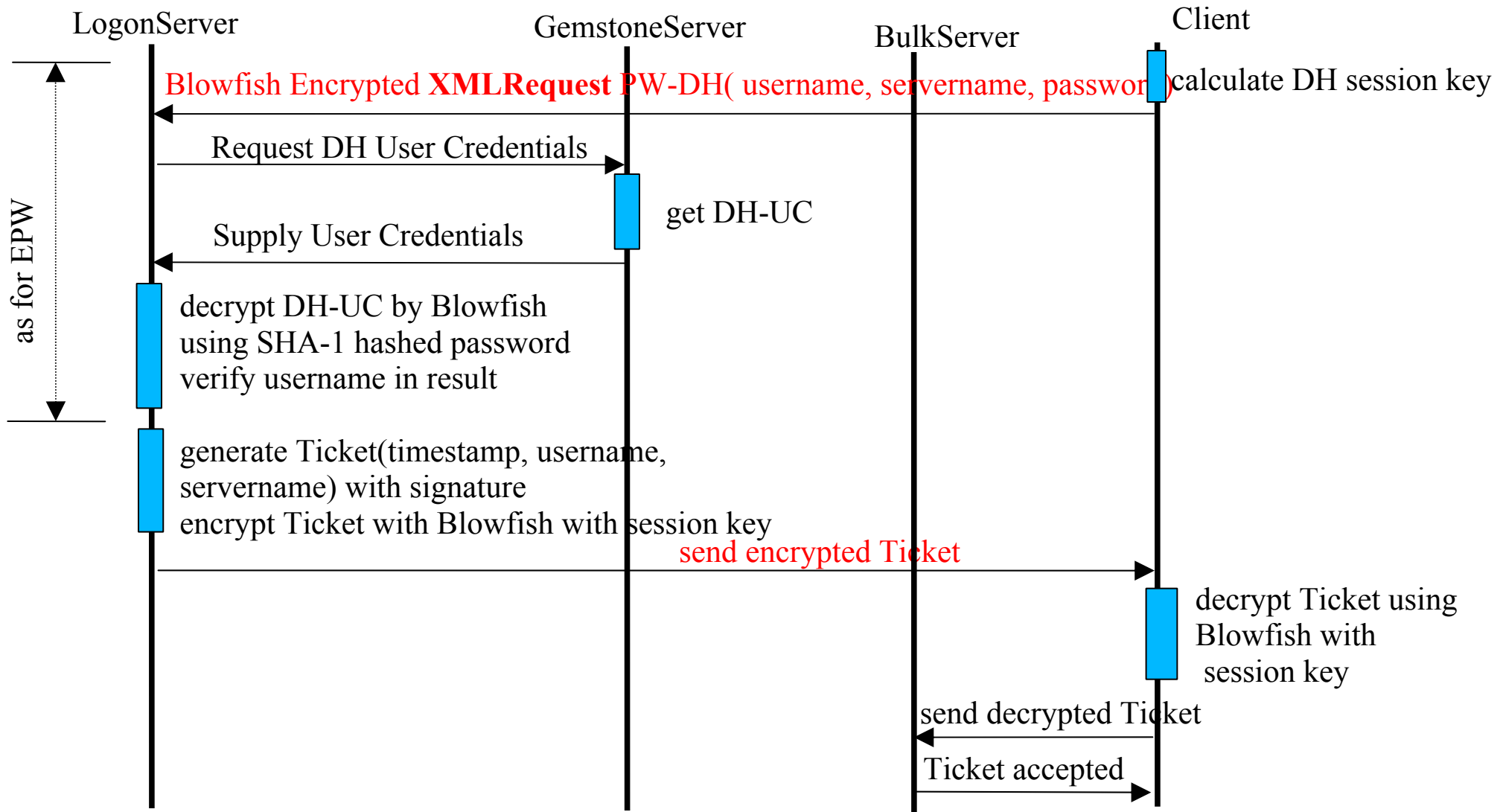
Tickets

# XML-Bulkserver

- OASE data can be requested from the so called XML-Bulkserver which is fully XML-based
  - Like LS a headless VW3.0-application
- The user must login into GS to access the data
- How can this be achieved when the user must supply an EPW?
  - by an XML-Ticket which allows access to GS

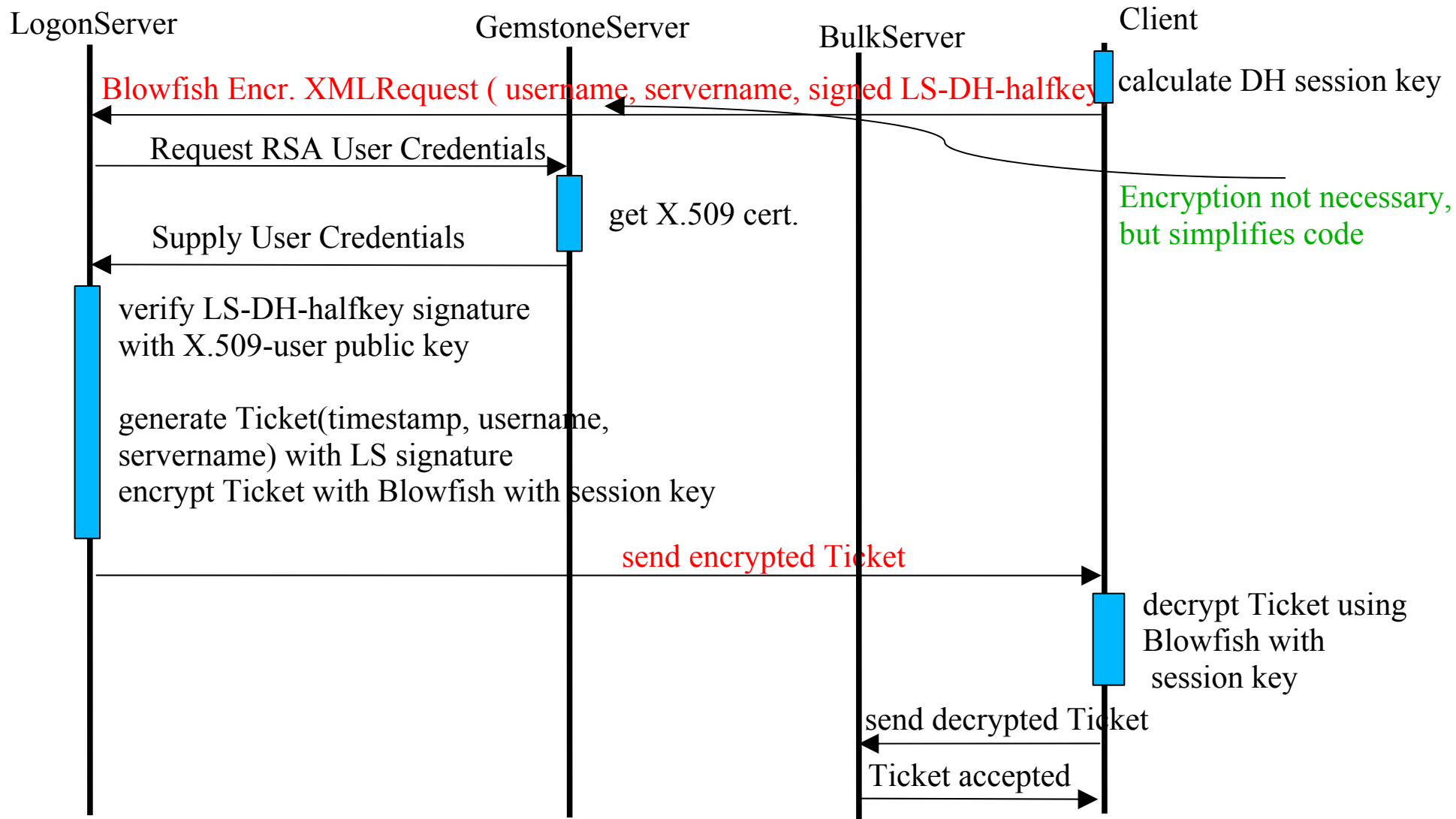


# XML-Ticket Overview for PW-DH





# XML-Ticket Overview for SC-RSA



# XML-Ticket basics

- The user requests an XML-Ticket from the LS
- When the user is allowed to access the database, LS creates a ticket
  - contains servername, username and a timestamp for validity checking.
  - the ticket is signed by LS
- The ticket must be presented to the Bulkserver within a short period (currently < 10 seconds).

# XML-Ticket checking

- The Bulkserver checks the ticket:
  - verifies the signature of LS
  - checks the timestamp
    - also verifies, that this ticket has not already been used
  - uses the servername to access GS
    - today, we do not restrict access for read for legal users, because of that there is no check necessary for the username

# XML-Ticket modes

- PW-DH and SC-RSA use the same technique for sending the ticket to the client:
  - the ticket format has no difference for both modes
  - the ticket is encrypted using blowfish with the DH-sessionkey

# XML-Ticket formats

- The TicketRequest differs for both modes
  - parameters: servername and username
- in SC-RSA the user signs the DH-halfkey from the LS with his signing key
- in PW-DH the credentials for servername, username are checked as for EPW using password
- these requests are both sent encrypted by blowfish

# TicketResponse

```
- <XMLTicket>  
  <Server>dpssM700</Server><User>awu</User>  
  <XMLTicketTime>3233767397000</XMLTicketTime>  
  <XMLTicketSignature>15022795967277155422365835660918884813150  
493043388515335841562848404653388507428348520470812988743204  
684100581555360979769072404615410077936309726151332490571843  
592914711119422903713287110926779070634207761829979664926899  
476112082718024066167433468885019163386939230270203100939218  
773849125730756896177136037</XMLTicketSignature>  
</XMLTicket>
```

- identical for both modes

# Encrypted TicketRequest Format

- both modes send encrypted with Blowfish:
  - `<XMLLogonRequest>`  
`<DhCode>4021923930769221418014505334054508087735352248076679`  
`0930681240902947674828299</DhCode>`  
`<EncryptedXMLLogonRequest>1013770570829590745205458228324819`  
`952813554860299057479356650031543473116011689133005574889073`  
`439715230885236966998448272871724104085738820437166598534732`  
`977536812429181457436359960820992836673414486482059363023010`  
`62332051528870082493408610500435018279009052905430115234</Enc`  
`ryptedXMLLogonRequest>`  
`</XMLLogonRequest>`

# Internal SC-RSA TicketRequest

```
- <XMLRSARequest>  
  <Server>dpssM700</Server><User>awu</User>  
  <Password>106942665803893961652237782571282335715752848048653  
488271645277633172859414869405045021075475426683763871005096  
512550742078122566663290545340367303214617163684950078288221  
884273972865236822265399295831080513274720629788515005292693  
213055680355212564287254503022928044905763429004378965569364  
748530842326832384</Password>  
</XMLRSARequest>
```

- <Password> contains the signature of the DH halfkey of the LS



# Internal PW-DH TicketRequest

- `<XmlEPWRequest>`  
  `<Server>dpssM700</Server><User>awu</User>`  
  `<Password>xxxxxxx</Password>`  
  `</XmlEPWRequest>`

- Same format as for EPW

# Batch-Tickets

- We plan, to introduce also Tickets for Batch-Processing
  - the user requests a ticket for a batch job at some time at some date
  - the LS delivers the ticket
    - the batch job is responsible for the secure deposit of the ticket until used
  - the batch job presents the decrypted ticket at the correct time/date and can be processed

# That's it!

## Questions?